

Simulating Floquet scrambling circuits on trapped-ion quantum computers

Kazuhiro Seki, Y.K., Tomoya Hayata, Seiji Yunoki, arxiv:2405.07613

Yuta Kikuchi (Quantinuum)

HHIQCD2024@YITP, Kyoto

Introduction

Quantum advantage for artificial tasks (RCS)

Google '19;
Zuchongzhi '21;
Quantinuum '24



Quantum dynamics?
(what else?)



Fault-tolerant QC

- Prime factoring
- QFT simulation
- Many more...



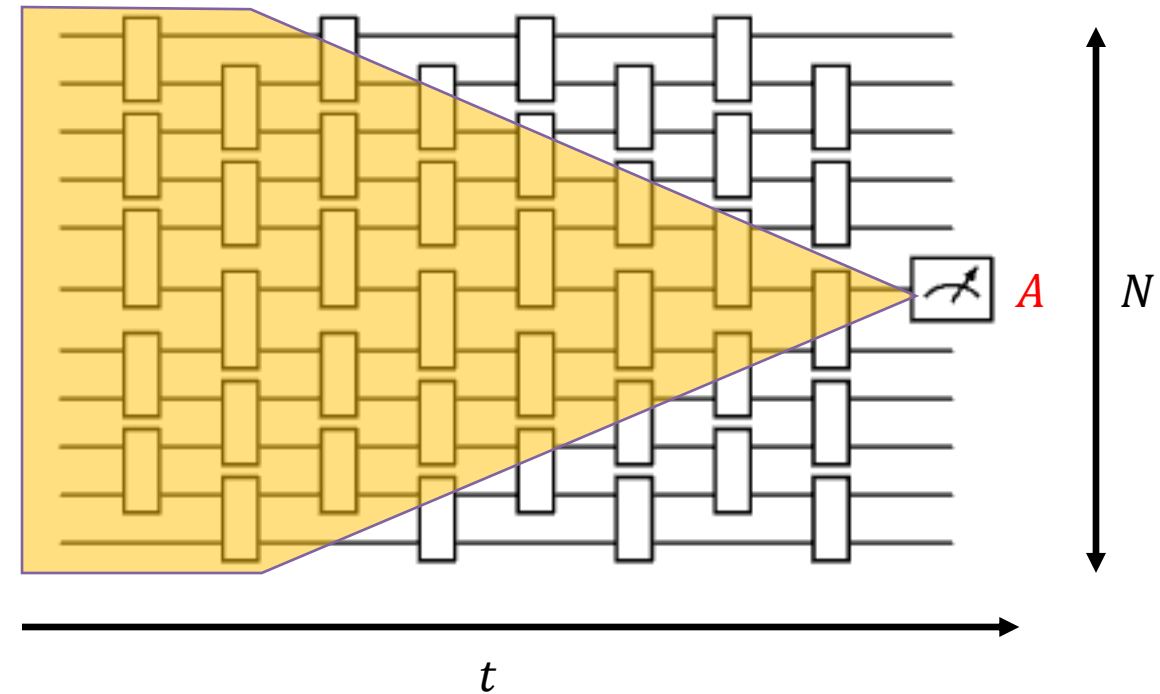
Need to cope with significant noise!!

(What can we do with **50-100 qubits** and **gate fidelity 99.9%, 99.99%, ... ?**)

Digital quantum simulation on noisy hardware

Setup

- Measure a local observable
- Circuit has geometrically local connectivity



Digital quantum simulation on noisy hardware

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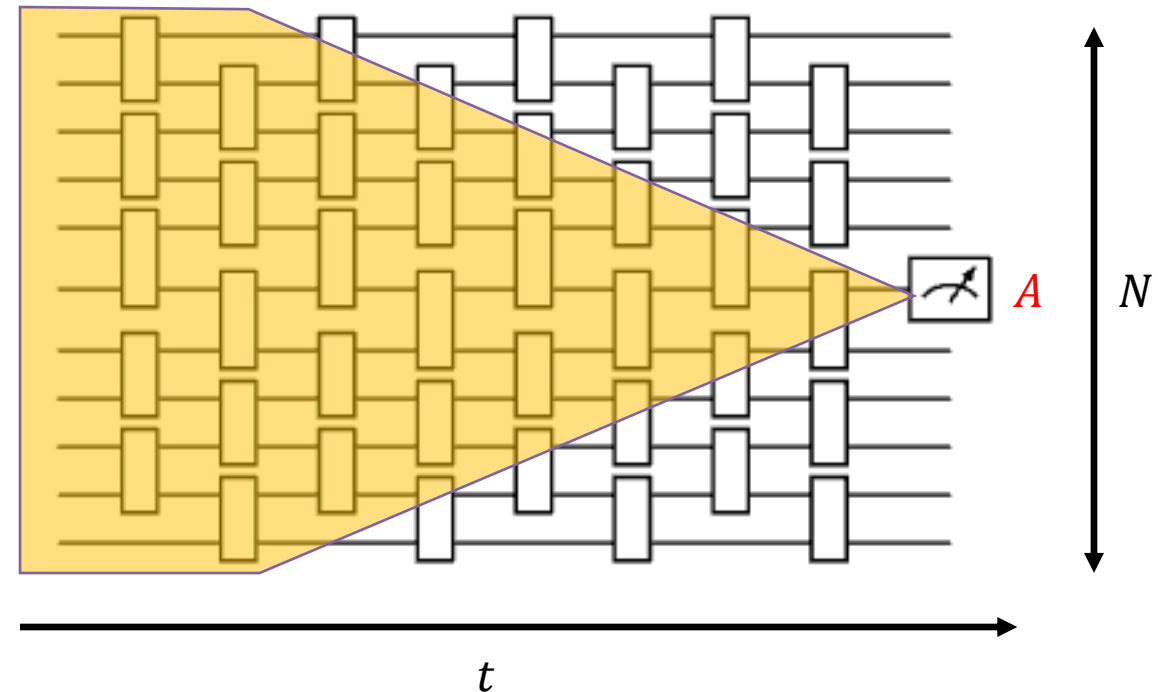
Crude estimate of error

- Dominant source of error is 2-qubit gates
- Yellow contains N_{2Q} 2-qubit gates
- Model the noisy circuit by the reduced density matrix

$$\rho_A^{\text{noisy}} = f \rho_A^{\text{ideal}} + (1 - f) \frac{I^{|A|}}{2^{|A|}}, \quad f = (1 - p_{2Q})^{N_{2Q}}$$

- Expectation value of traceless operator A is

$$\text{Tr}_A[A \rho_A^{\text{noisy}}] = f \text{Tr}_A[A \rho_A^{\text{ideal}}]$$



Digital quantum simulation on noisy hardware

Hamiltonian simulation by Trotterization

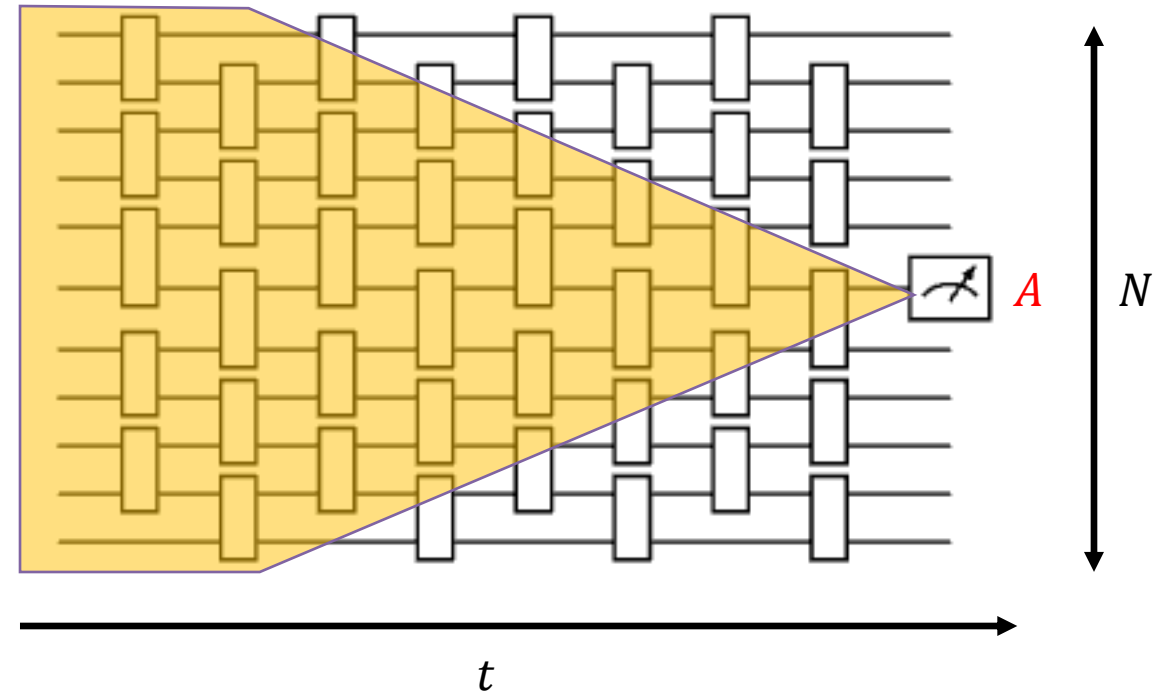
$$H = H_1 + H_2$$

$$e^{-iHt} \approx \left(e^{-iH_1T} e^{-iH_2T} \right)^{t/T}$$

Crude estimate of error

- Parameters: $t = 5$, $T = 0.1$, $N = 20$
- # of 2Q gates: $N_{2Q} \approx N \frac{t}{T} = 1000$
- 2Q gate error: $p_{2Q} = 2 \times 10^{-3}$
- Expectation value of traceless operator A is

$$\text{Tr}_A[A\rho_A^{\text{noisy}}] = f \text{Tr}_A[A\rho_A^{\text{ideal}}], \quad f = (1 - p_{2Q})^{N_{2Q}} = 0.135$$



Floquet dynamics

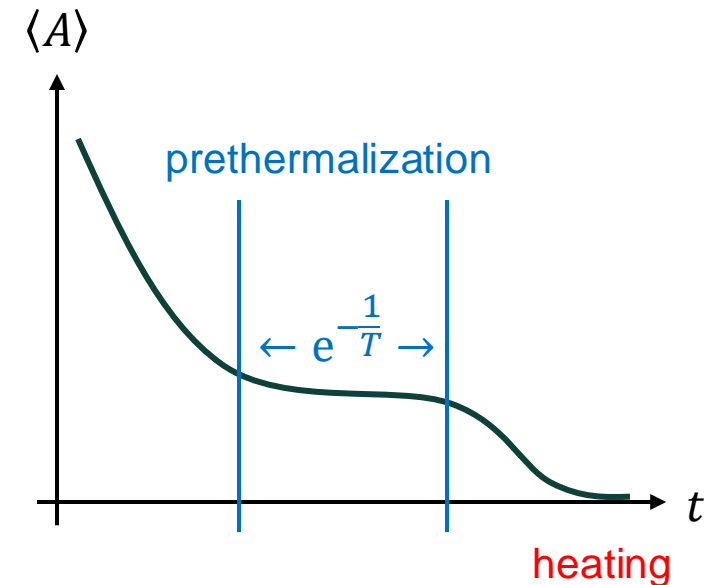
Floquet dynamics:

- Periodically driven dynamics described by Hamiltonian $H(t + T) = H(t)$
- Floquet systems eventually heat up to infinite temperature by acquiring energy from the driving force

(Floquet heating)

$$\langle \psi(t) | A | \psi(t) \rangle \approx \text{Tr}[A(I^N / 2^N)] \quad \text{for} \quad |\psi(t)\rangle = \mathcal{T}e^{-i \int_0^t dt' H(t')} |\psi(0)\rangle$$

Lanzarides, Das & Moessner '14;
D'Alessio & Rigol '14;
Abanin, Roeck & Huveneers '15;
Mori, Kuwahara & Saito '16



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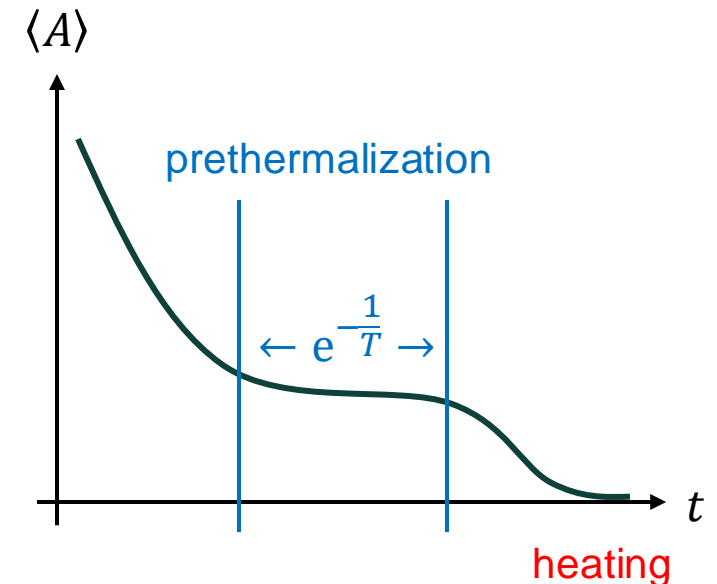
Lazarides, Das & Moessner '14;
D'Alessio & Rigol '14;
Abanin, Roeck & Huvneers '15;
Mori, Kuwahara & Saito '16

Trotterization

- Trotter dynamics for time mT , $(e^{-iH_1 T/2} e^{-iH_2 T/2})^{2m}$, is a Floquet dynamics of m cycles
- Qualitative change happens as the period T is increased

→ Floquet heating happens during Trotter dynamics

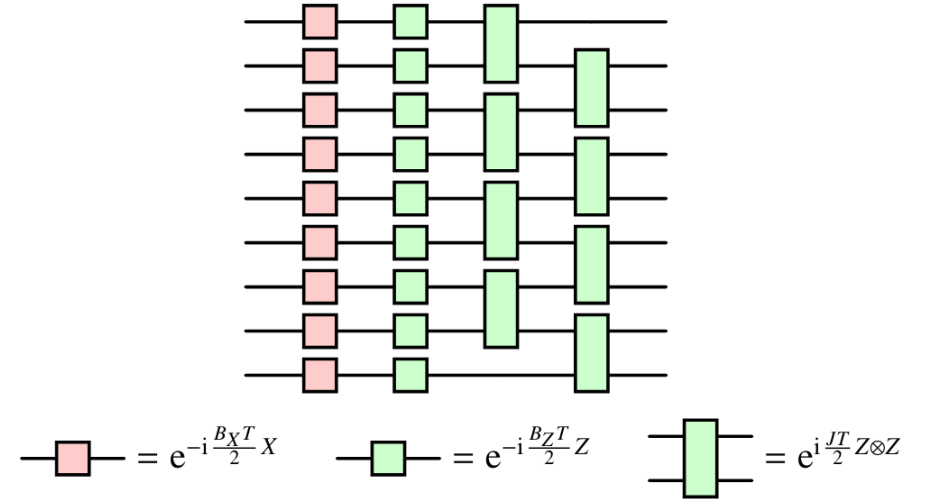
Heyl, Hauke & Zoller '18;
Varnier, Bertini, Giudici & Piroli '23



Information scrambling by Floquet circuits

Kicked-Ising model Prosen '02, '07

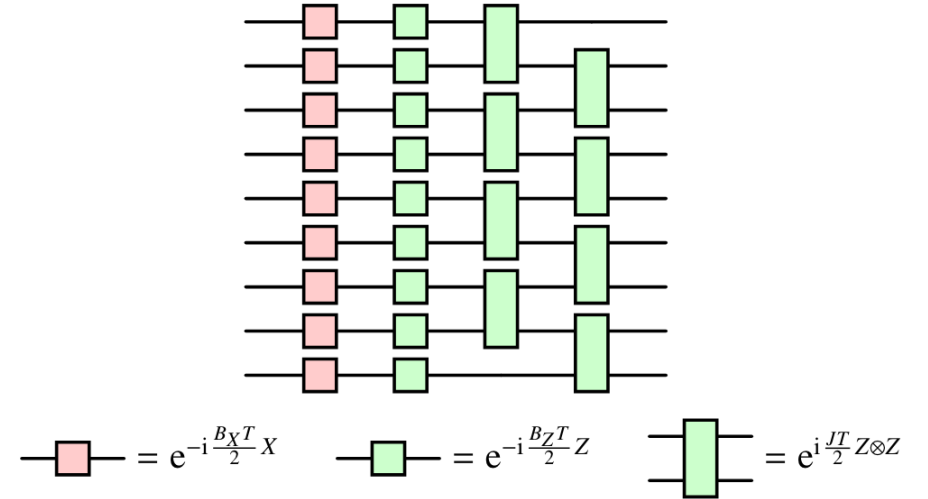
- $U_F = e^{-iH_Z \frac{T}{2}} e^{-iH_X \frac{T}{2}},$
- $H_Z = -J \sum_i Z_i Z_{i+1} + B_Z \sum_i Z_i, \quad H_X = B_X \sum_i X_i$
- **Maximally chaotic point at $|JT| = |B_X T| = \frac{\pi}{2}$**



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Information scrambling:

- A process of the lost information spreading across the system
 - Scrambling dynamics makes it hard to recover the initial information
- **Diagnose the complexity of dynamics**

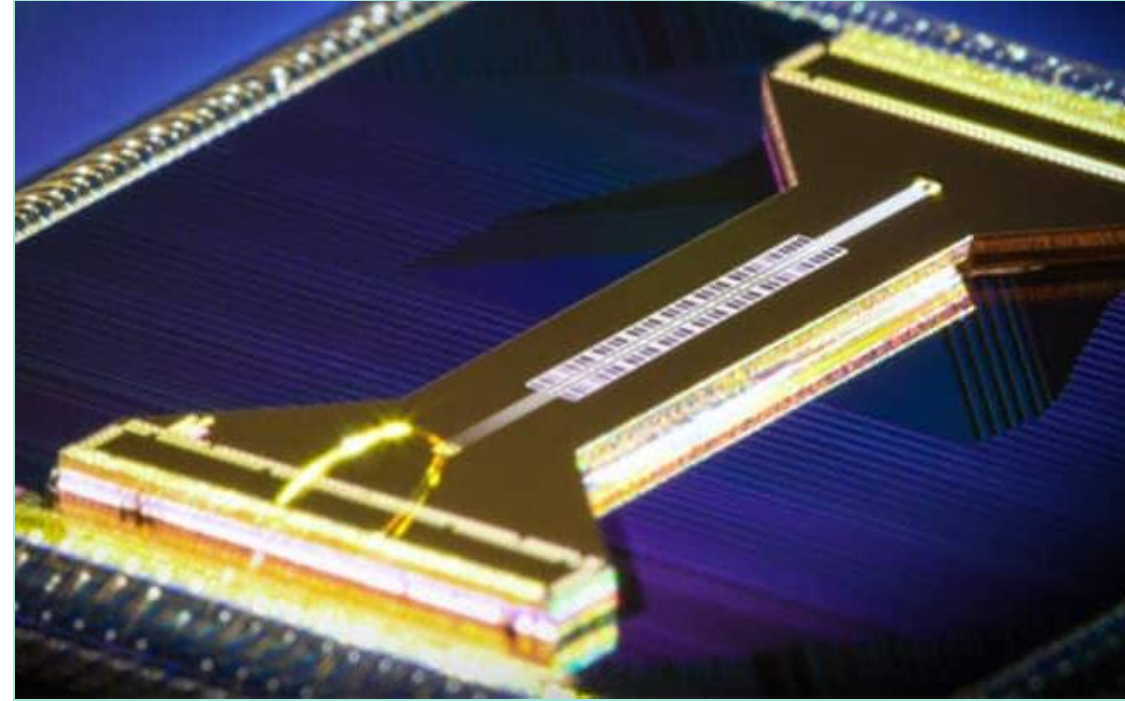
Experiments on trapped-ion quantum computers

GOAL:

Access the feasibility of scrambling simulation on the current hardware

Seki, YK, Hayata, Yunoki '24

- Hayden-Preskill recovery protocol
- Interferometric protocol for OTOCs
- Thermal expectation value with microcanonical TPQ states



- 99.91% Two-qubit gate fidelity (arbitrary angle)
- 20 qubits
- 99.998% Single-qubit gate fidelity
- Measurement cross talk error < 0.01%
- All-to-all-connectivity
- SPAM fidelity > 99.7%

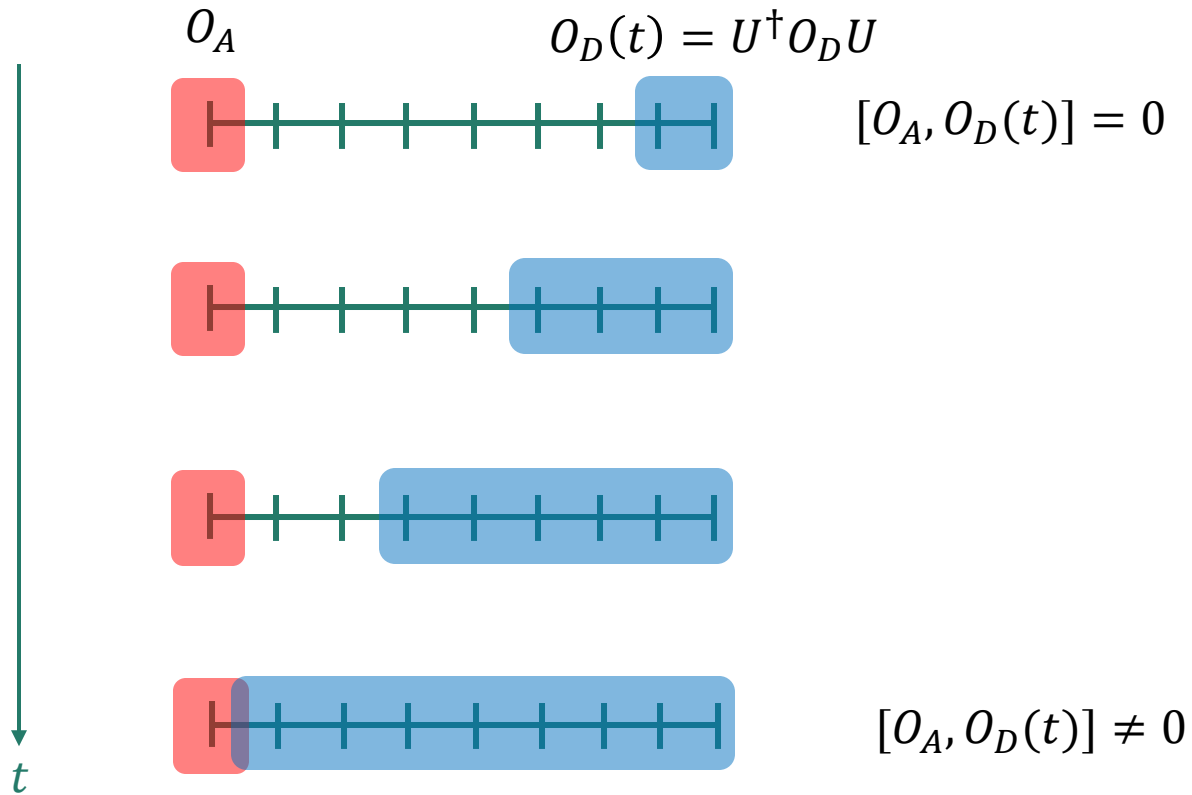
System Model H1
Available on various platforms



OTOCs

Larkin & Ovchinnikov 1969;
Shenker & Stanford 2014

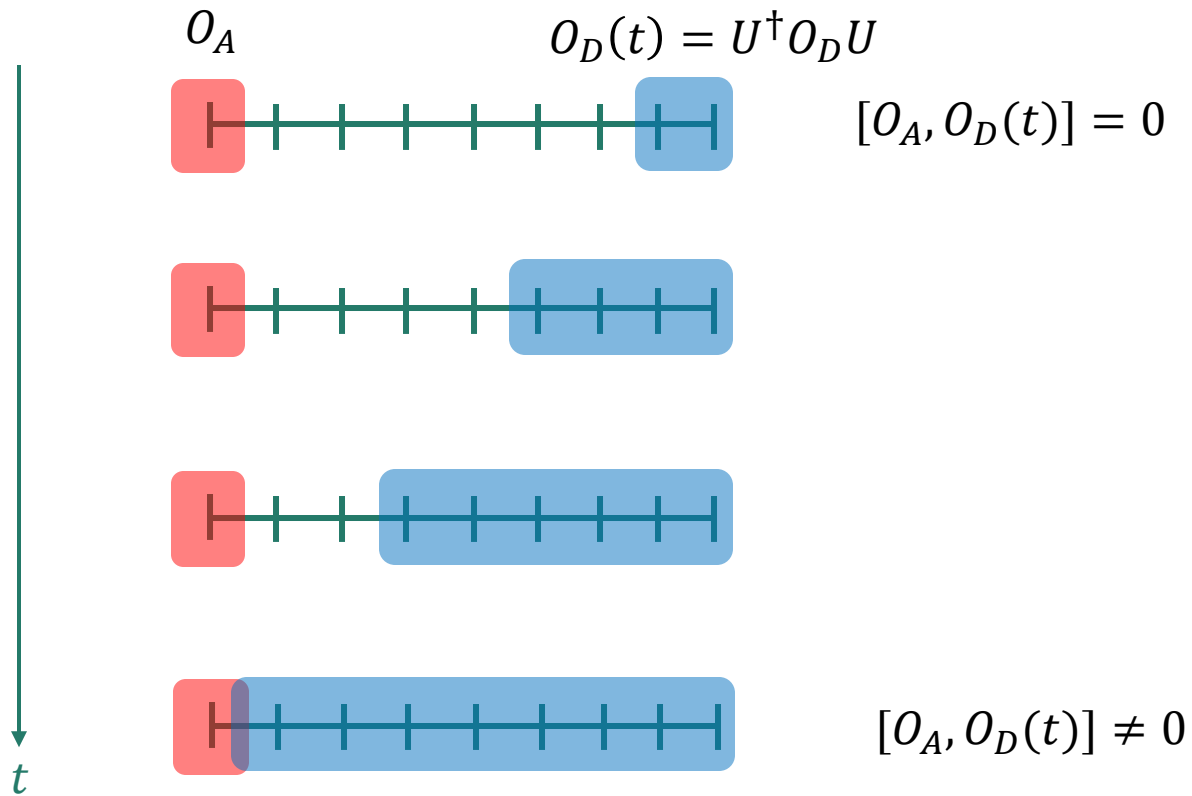
Operator growth



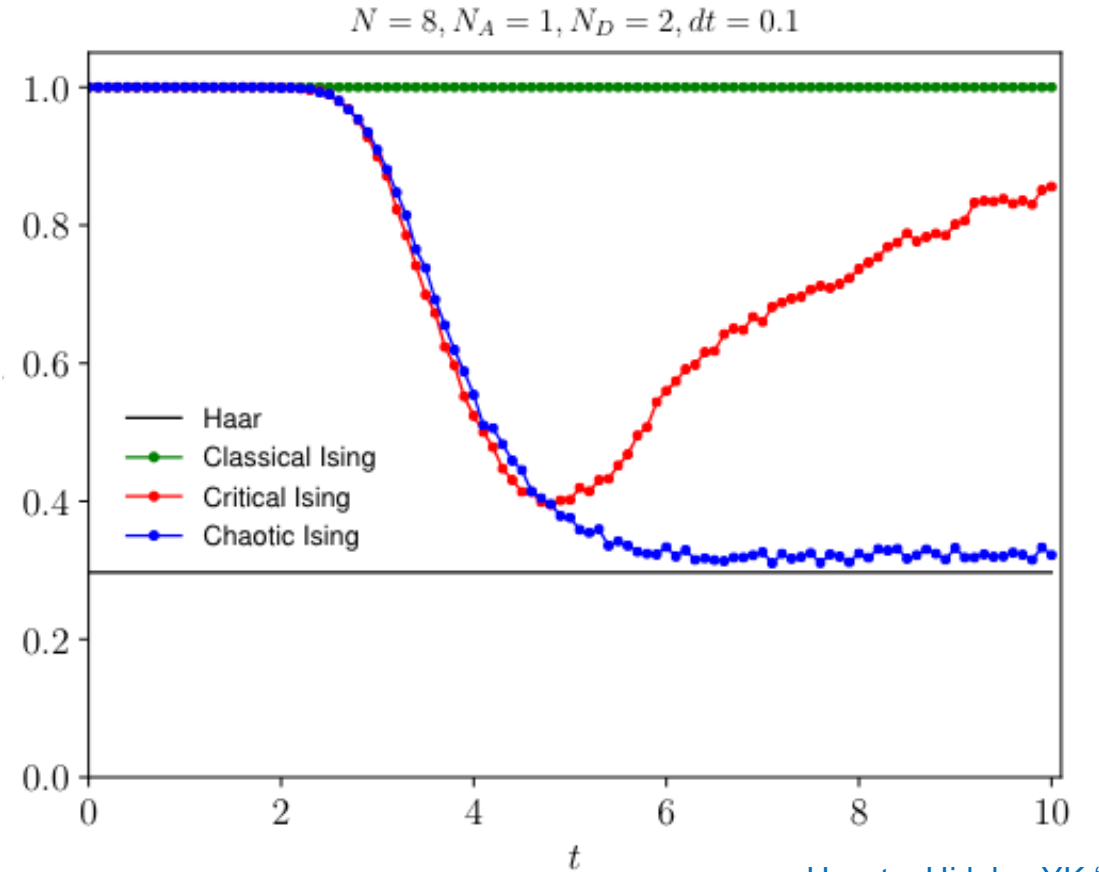
OTOCs

Larkin & Ovchinnikov 1969;
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Operator growth



$$\text{OTOC} = \langle O_A O_D(t) O_A^\dagger O_D(t)^\dagger \rangle$$



Hayata, Hidaka, YK '21

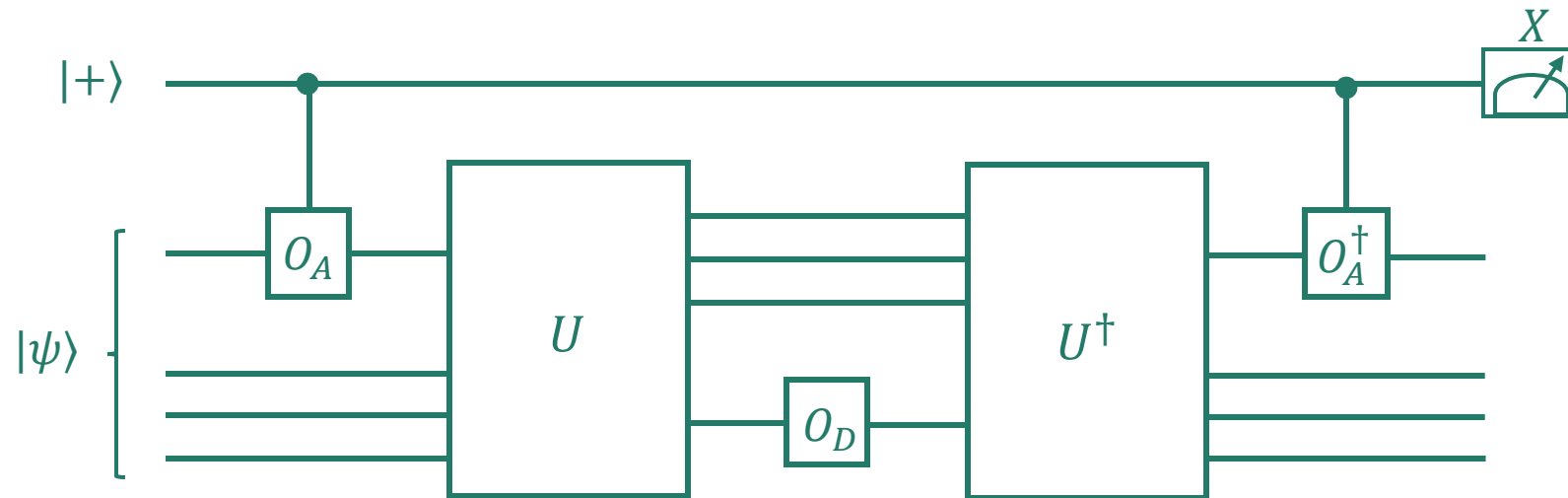
$$H = -J \sum_i Z_i Z_{i+1} + B_Z \sum_i Z_i + B_X \sum_i X_i$$



OTOCs

Interferometric protocol

$$\text{OTOC} = \langle O_A O_D(t) O_A^\dagger O_D(t)^\dagger \rangle$$



Swingle, Bentsen, Schleier-Smith & Hayden '16
Swingle & Halpern '18
Mi et.al. '21

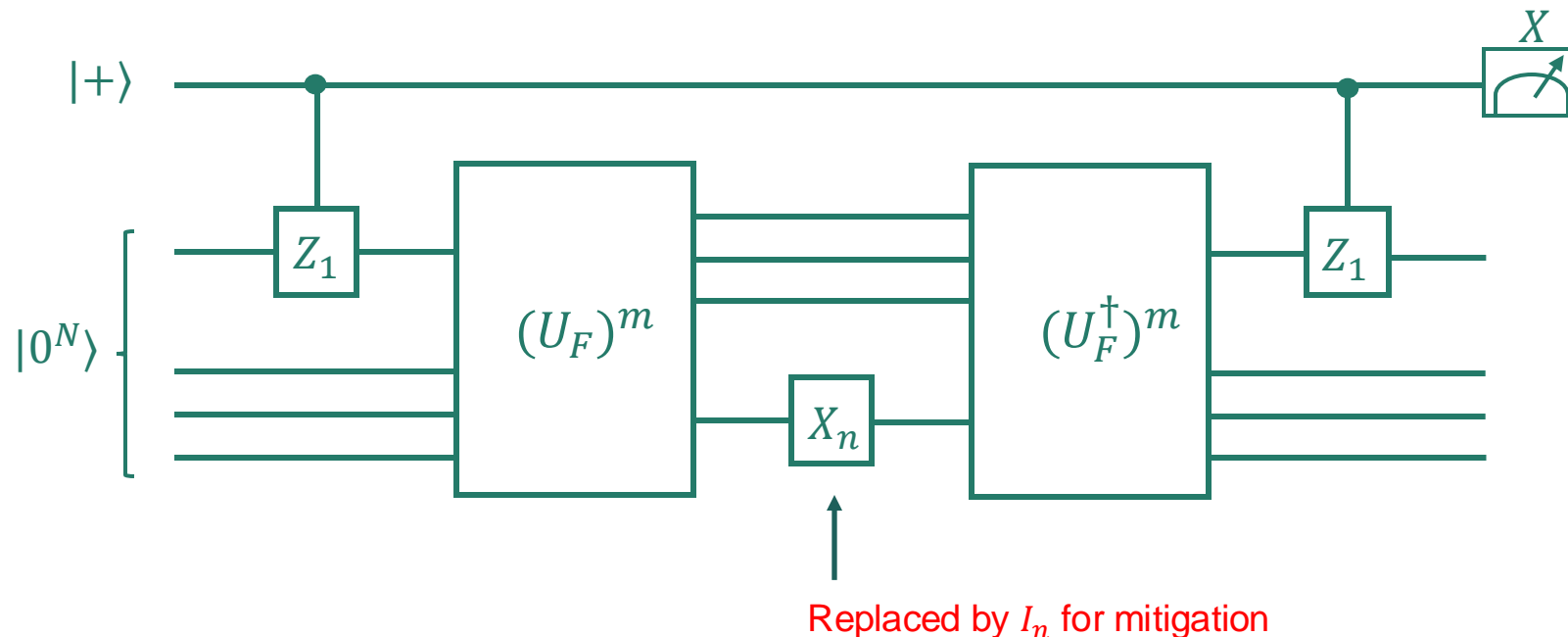
OTOCs

Interferometric protocol

$$\text{OTOC} = \langle Z_1 X_n(t) Z_1 X_n(t) \rangle$$

$$\text{OTOC}' = \frac{\langle Z_1 X_n(t) Z_1 X_n(t) \rangle}{\langle Z_1 I_n(t) Z_1 I_n(t) \rangle} : \text{mitigates incoherent errors}$$

$$\langle Z_1 I_n(t) Z_1 I_n(t) \rangle = 1 \text{ w/o noise}$$



Swingle, Bentsen, Schleier-Smith & Hayden '16
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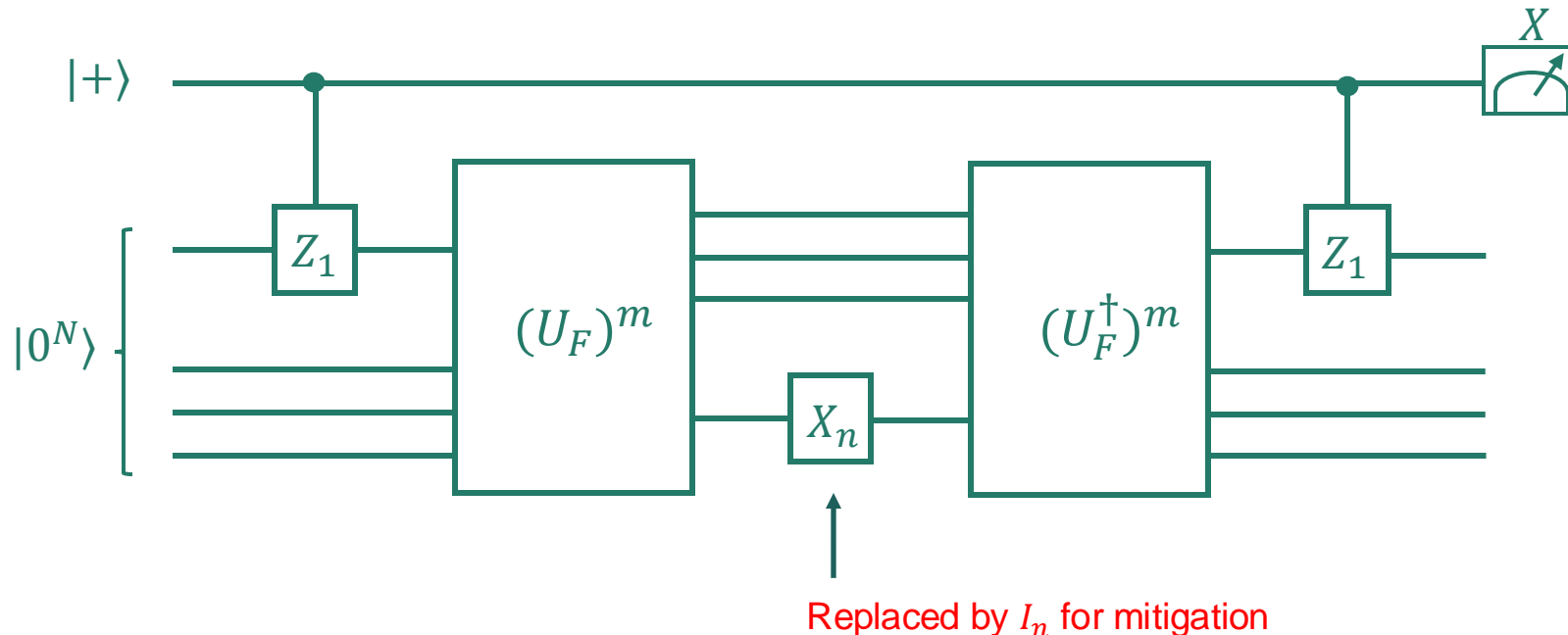
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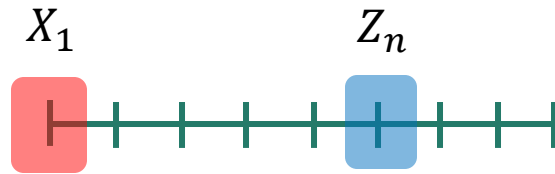
Setup

- 19-qubit spin chain + 1 ancilla qubit
- $U_F = e^{-iH_Z \frac{T}{2}} e^{-iH_X \frac{T}{2}}$
- 367 two-qubit gates inside the causal cone at maximum

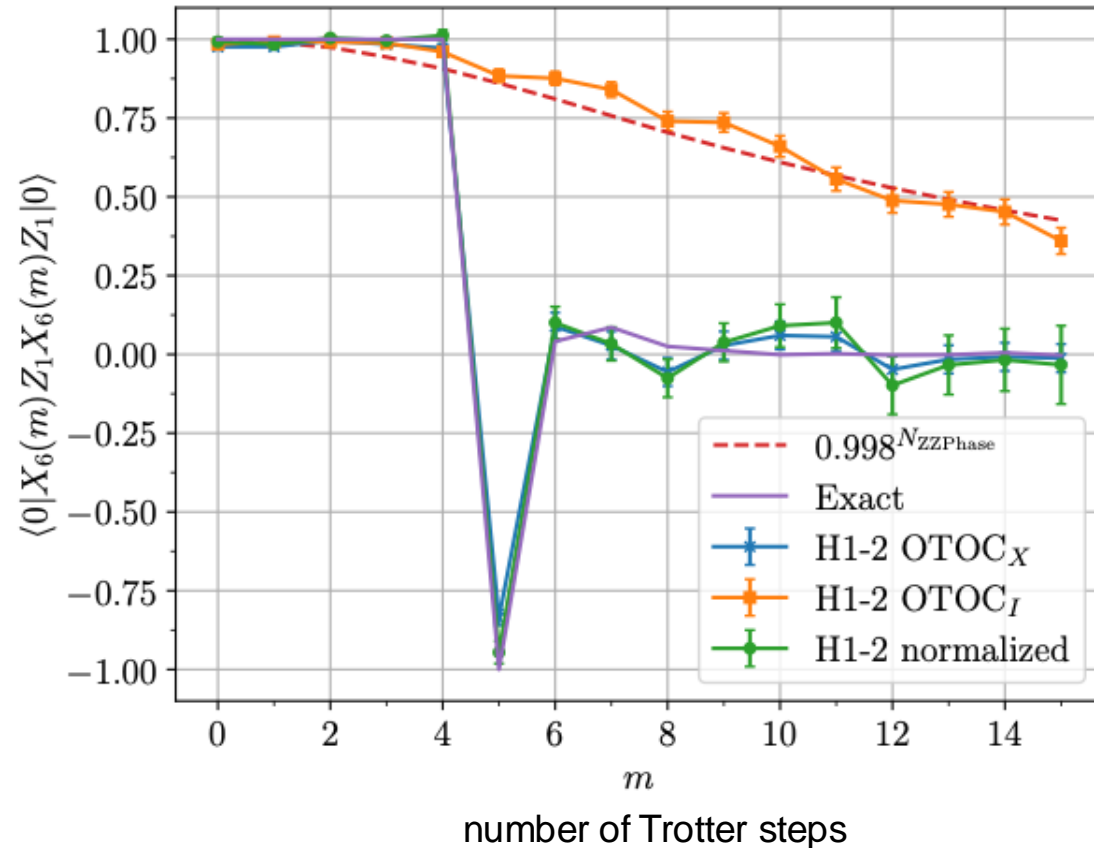
Swingle, Bentsen, Schleier-Smith & Hayden '16
Swingle & Halpern '18
Mi et.al. '21

OTOCs [Experiment]

Seki, YK, Hayata, Yunoki '24



$n = 6$



$$\text{OTOC}_X = \langle Z_1 X_n(t) Z_1 X_n(t) \rangle$$

$$\text{OTOC}_I = \langle Z_1 I_n(t) Z_1 I_n(t) \rangle \approx (1 - p_{2Q})^{N_{2Q}} \approx 0.998^{N_{2Q}}$$

Observations

- OTOC_I decays due to hardware noise
- $(1 - p_{2Q})^{N_{2Q}}$ approximates OTOC_I very well

OTOCs [Experiment]

Seki, YK, Hayata, Yunoki '24

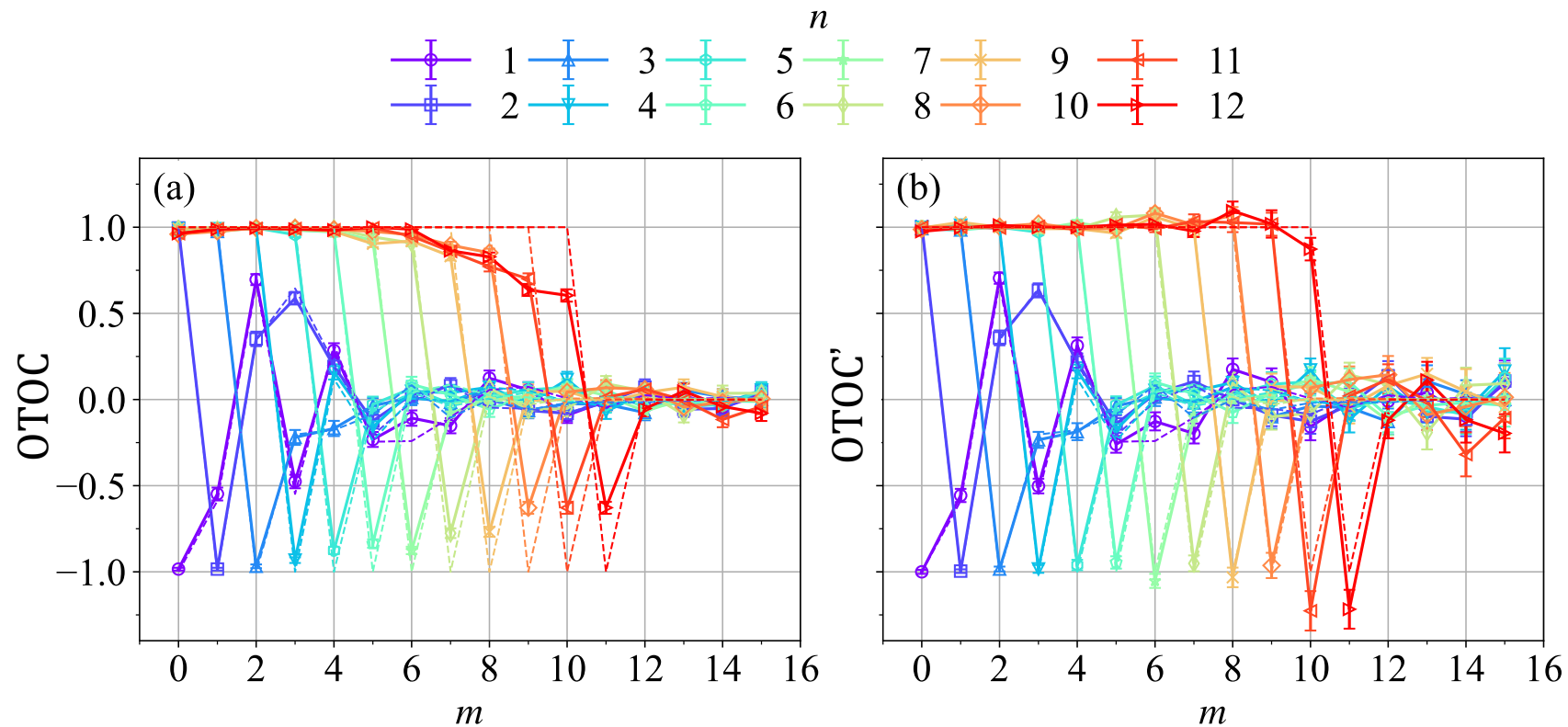
Dashed: exact

Left: $OTOC = \langle Z_1 X_n(t) Z_1 X_n(t) \rangle$

Right: $OTOC' = \frac{\langle Z_1 X_n(t) Z_1 X_n(t) \rangle}{\langle Z_1 I_n(t) Z_1 I_n(t) \rangle}$

Observations

- $OTOC'$ suffers less from noise at early times
- Statistical error of $OTOC'$ is amplified at late times



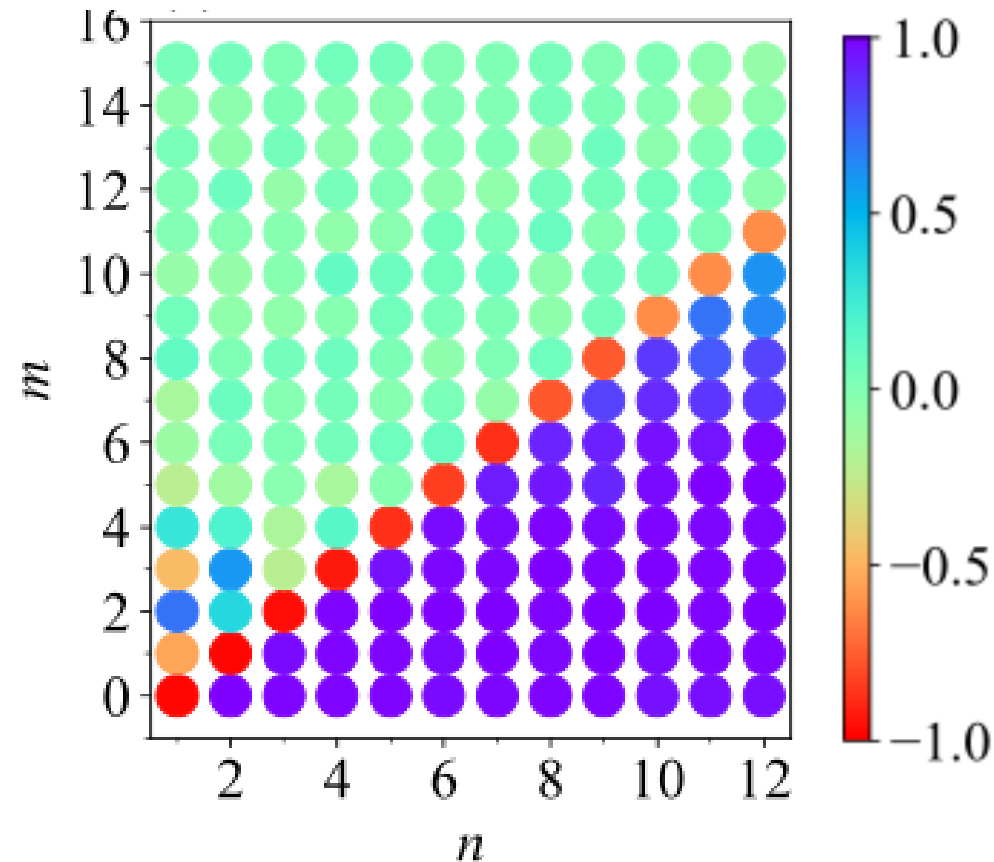
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Seki, YK, Hayata, Yunoki '24

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Observations

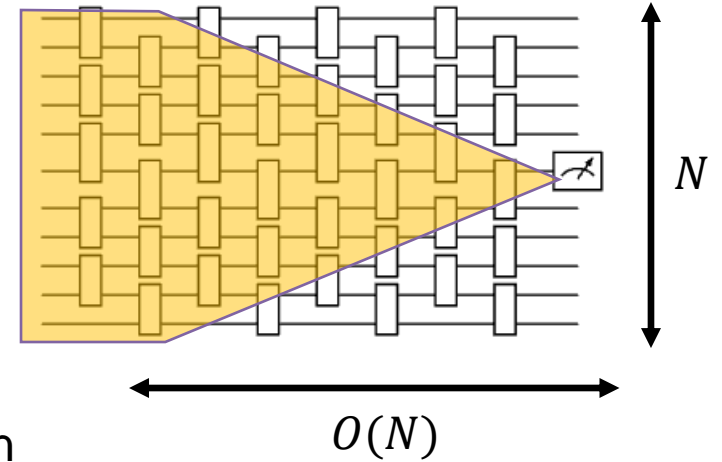
- Ballistic growth of entanglement
- Statistical error of OTOC' is amplified at late times



Discussion

Circuit fidelity is crudely (sometimes accurately) estimated by

$$f = (1 - p_{2Q})^{N_{2Q}}$$



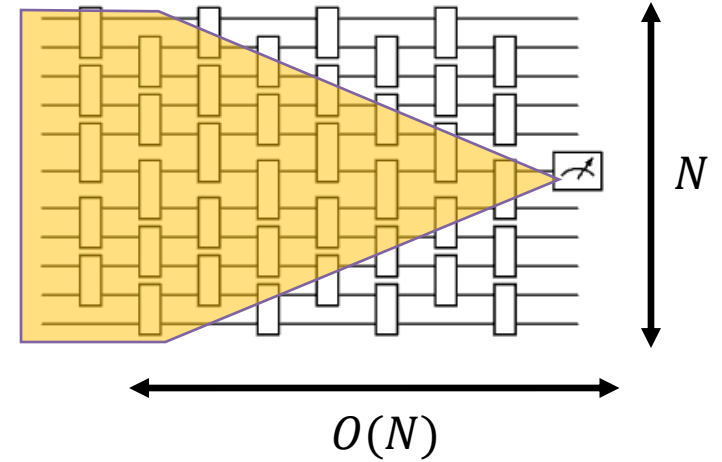
Geometrically local models requires $\text{poly}(N)$ gates to scramble information

- 1D geometrically circuit requires $t \approx O(N)$ for the entire system to get involved: $N_{\text{gate}} \sim O(N^2)$
- E.g. we used 400 2Q gates for the 20-qubit 1D system
 \Rightarrow a 40-qubit system requires 4 times larger gate counts ~ 1600 2Q gates : $f = 0.999^{1600} \approx 0.2$

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What if hardware result deviates from the estimate?

- Memory error (on idling qubits, during ion shuttling)
- SPAM error (bias between $0 \rightarrow 1$ and $1 \rightarrow 0$)
- Gate counting analysis overestimates the error

Schiffer, Rubio Trivedi & Cirac '24
Granet & Dreyer '24
Chertkov, Chen, Lubasch, Hayes, Foss-Feig '24

